

# Counting & Generating Irreducible Quadratics over $\mathbb{Z}_n[x]$

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## Introduction

It is known that there are  $\binom{p^k}{2}$  irreducible, quadratic, monic polynomials in  $\text{GF}(p^k)[x]$  where  $p$  is a prime and  $k \in \mathbb{N}$ . My research began with two questions.

- How many are there in  $\mathbb{Z}_n[x]$  for any integer  $n$ ?
- How can we efficiently generate these polynomials?

### Example: Quadratics over $\mathbb{Z}_4$

	Reducible	Irreducible
$x^2$	$= (x+0)^2 = (x+2)^2$	$x^2 + 1$
$x^2 + x$	$= (x+0)(x+1)$	$x^2 + 3$
$x^2 + 2x$	$= (x+0)(x+2)$	$x^2 + x + 1$
$x^2 + 3x$	$= (x+0)(x+3)$	$x^2 + x + 3$
$x^2 + 2x + 1$	$= (x+1)^2 = (x+3)^2$	$x^2 + 2x + 2$
$x^2 + 3x + 2$	$= (x+1)(x+2)$	$x^2 + 2x + 3$
$x^2 + 3$	$= (x+1)(x+3)$	$x^2 + 3x + 1$
$x^2 + x + 2$	$= (x+2)(x+3)$	$x^2 + 3x + 3$

We may have expected to find  $\binom{2^2}{2} = 6$  irreducible quadratics, but in fact there are 8.

## Computational Experiments

Let  $C(n)$  be the number of irreducible, quadratic, monic polynomials in  $\mathbb{Z}_n[x]$ .

Using Algorithm 1 coded in C, we found the following.

### Sample Data from Algorithm 1

$i$	$C(2^i)$	$C(3^i)$	$C(5^i)$	$C(7^i)$
1	1	3	10	21
2	8	45	350	1323
3	36	432	9000	65856
4	160	4050	227500	3241350
5	656	36693	5693750	158876571

### Conjecture

Let  $p$  be an odd prime and  $k \in \mathbb{N}$ . Then

$$C(2^k) = \begin{cases} \frac{1}{3}2^{2k+1} - \frac{4}{6}2^k & \text{if } k \text{ is even} \\ \frac{1}{3}2^{2k+1} - \frac{5}{6}2^k & \text{if } k \text{ is odd} \end{cases}$$

$$C(p^k) = \begin{cases} p^{2k} - \frac{p^k(p^{k+1}+p+2)}{2(p+1)} & \text{if } k \text{ is even} \\ p^{2k} - \frac{p^k(p^{k+1}+2p+1)}{2(p+1)} & \text{if } k \text{ is odd} \end{cases}$$

Next we wanted to prove this conjecture, determine a general formula, and write an efficient algorithm to generate irreducible quadratics.

## Proof Approach

Consider the quadratic, monic polynomials in  $\mathbb{Z}_n[x]$ . We have

$$n^2 = C(n) + R(n)$$

where  $R(n)$  is the number of reducible.

Then  $R(n)$  is equal to the number of integer pairs  $(b, c)$  such that  $0 \leq b, c < n$  and

$$x^2 + bx + c \equiv 0 \pmod{n}$$

has a solution.

From Nagell, we see that this congruence has a solution if and only if the following system of congruences, has a solution.

$$y^2 \equiv b^2 - 4c \pmod{4n} \quad (1)$$

$$y \equiv b \pmod{2} \quad (2)$$

Congruence (1) suggests that solving for  $R(n)$  and  $C(n)$  begins by counting squares in  $\mathbb{Z}_{4n}$

## Counting $R(n)$

We use the following 3 observations from Stangl to count squares in  $\mathbb{Z}_{p^i}$  where  $p$  is a prime and  $i \geq 3$  an integer.

- Let  $S(n)$  is the number of squares in  $\mathbb{Z}_n$  and  $Q(n)$  the number of quadratic residues in  $\mathbb{Z}_n$ .

$$S(p^i) = Q(p^i) + S(p^{i-2})$$

- If  $p$  is an odd prime, then

$$Q(p^i) = \frac{\phi(p^i)}{2} = \frac{p^i - p^{i-1}}{2}$$

- For  $p = 2$  we have

$$Q(2^i) = \frac{2^{i-1}}{4} = 2^{i-3}$$

Finally, by the previous congruences and the Chinese Remainder Theorem,

$$R(n) = \frac{n}{2} S(2^{\alpha_0+2}) \prod_{i=1}^k S(p_i^{\alpha_i})$$

when  $n$  has prime factorization

$$n = p_0^{\alpha_0} \times p_1^{\alpha_1} \times \dots \times p_k^{\alpha_k}$$

## Main Theorem

Suppose  $n$  has prime factorization

$$n = 2^{\alpha_0} \times p_1^{\alpha_1} \times p_2^{\alpha_2} \times \dots \times p_k^{\alpha_k}$$

where  $p_i$  are odd primes and  $\alpha_i$  are non-negative integers.

Then the number of irreducible, quadratic, monic polynomials in  $\mathbb{Z}_n[x]$  is

$$C(n) = n^2 - \frac{n}{2} S(2^{\alpha_0+2}) \prod_{i=1}^k S(p_i^{\alpha_i})$$

where

$$S(p^i) = \begin{cases} \frac{p^{i+1}+p+2}{2(p+1)} & \text{if } i \text{ is even} \\ \frac{p^{i+1}+2p+1}{2(p+1)} & \text{if } i \text{ is odd} \end{cases} \quad S(2^i) = \begin{cases} \frac{2^{i-1}+4}{3} & \text{if } i \text{ is even} \\ \frac{2^{i-1}+5}{3} & \text{if } i \text{ is odd} \end{cases}$$

when  $p$  is an odd prime.

## References

Stangl, W. D. (1996). Counting Squares in  $\mathbb{Z}_n$ . *Mathematics Magazine*, 69(4), pp. 285-289. <https://doi.org/10.1080/0025570X.1996.11996456>

Nagell, T. (1964). *Introduction to Number Theory* (2nd ed.). New York: Chelsea Publishing Company.

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## Generating Reducible Quadratics

### Algorithm 1

```

1: P ← []; i ← 0
2: for b ∈ Z_{p^k} do
3:   for c ∈ Z_{p^k} do
4:     for x_0 ∈ Z_{p^k} do
5:       if x_0^2 + bx_0 + c ≡ 0 (mod p^k) then
6:         P[i] ← {x^2 + bx + c}; i++
7:       break
8:     end if
9:   end for
10: end for
11: end for
12: return P

```

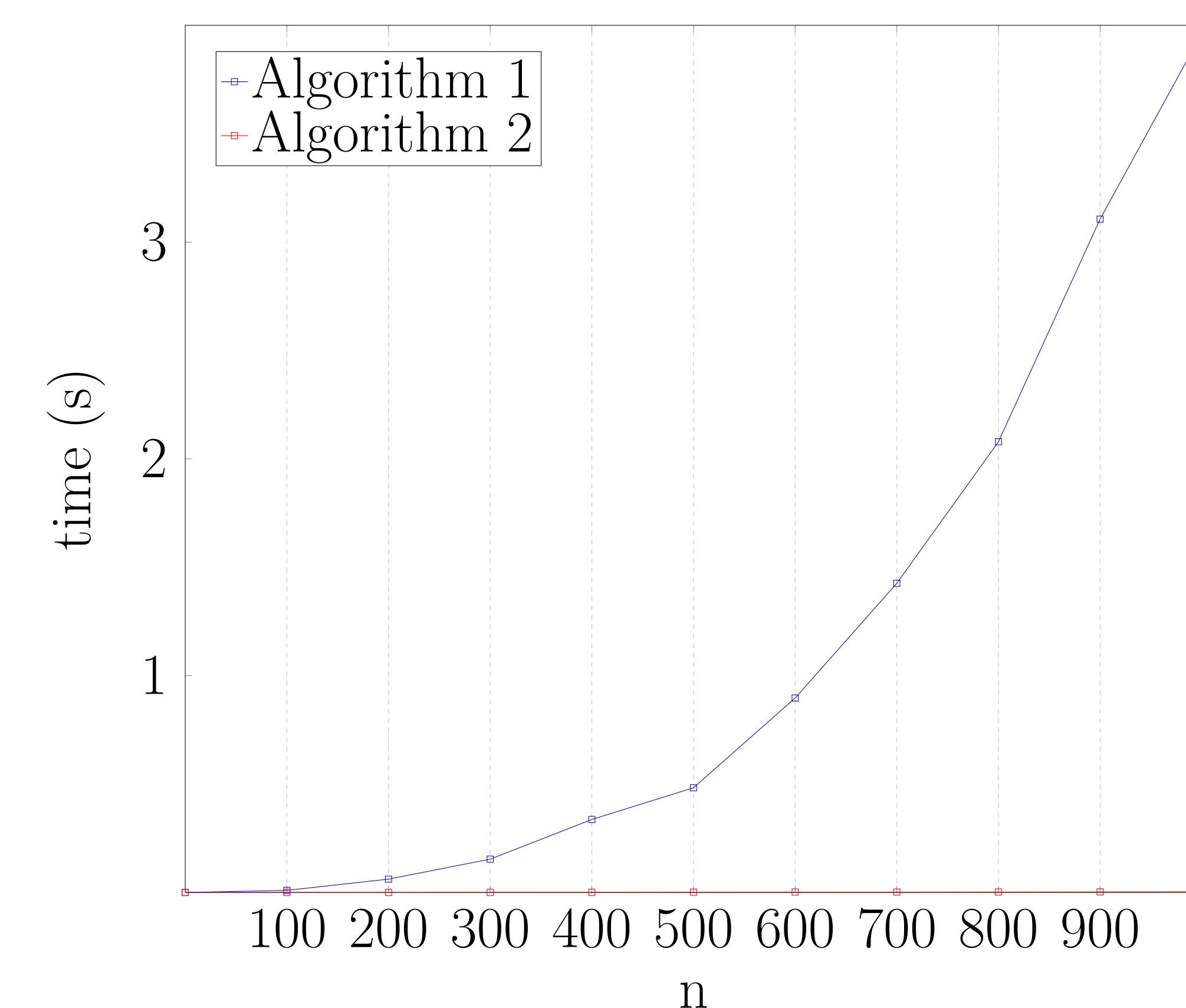
### Algorithm 2

```

1: P ← []; i ← 0
2: for s ∈ {x^2 mod 4n : x ∈ Z_{4n}} do
3:   if 2|s then
4:     B ← {b ∈ Z_n : 2 | b}
5:   else
6:     B ← {b ∈ Z_n : 2 ∤ b}
7:   end if
8:   for b ∈ B do
9:     c ← (b^2 - s) / 4 mod n
10:    P[i] ← {x^2 + bx + c}; i++
11:  end for
12: end for
13: return P

```

Implemented in C



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